

# Nico's Numerical Nightmare: An upper bound on $p$ for disconnected $G_p$

Nico Tripeny

August 11, 2021

# The Cage

## Definition (The Cage)

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## Theorem

*If  $p$  is greater than  $10^{1310}$ , then  $G_p$  is connected.*

# Order of an Element

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The *order of a triple*  $(x, y, z)$  is the maximum of the order of its elements  $x$ ,  $y$ , and  $z$ .

## Order of an Element (Continued)

### Lemma

*If  $x = h + h^{-1}$  with  $h \in \mathbb{F}_p$  or  $h \in \mathbb{F}_{p^2}$ , then the order of  $x$  is the order of  $h$  in  $\mathbb{F}_p^*$  or  $\mathbb{F}_{p^2}^*$  (BGS)*



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### Example

$h \in \mathbb{F}_p$  Let  $p = 7$  and  $x = 6$ . We want to find  $h$  such that  $h + h^{-1} = 6$ .  
Subtracting 4 and multiplying by  $h$  we get  $h^2 - 6h + 1 = 0$

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$$h = \frac{6 \pm \sqrt{36-4}}{2} = \frac{6 \pm 2}{2} = 4, 2.$$

Since  $4^{-1} = 2$  in  $\mathbb{F}_p$ , we see this is the correct solution for  $h$ .

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Similar to above we get  $h = \frac{4 \pm \sqrt{16-4}}{2} = \frac{4 \pm \sqrt{5}}{2}$ . Since 5 is not a quadratic residue mod 7, this is in  $\mathbb{F}_{p^2}$  but is still a solution to  $h + h^{-1} = 4$

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### Lemma

*If  $x = h + h^{-1}$  with  $h \in \mathbb{F}_{p^2}$  then the order of  $h$  divides  $p + 1$  or  $p - 1$*

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### Proof.

Since  $x \in \mathbb{F}_p^*$ ,  $x^{p-1} = 1$ , we have  $x^p = (h + h^{-1})^p = h^p + h^{-p} = x$ .

Since  $h$  and  $h^{-1}$  are the only solutions to  $x = y + y^{-1}$ , we see  $h^p = h$  or  $h^p = h^{-1}$ .

So,  $h^{p+1} = 1$  or  $h^{p-1} = 1$  giving the desired result. □



# Orbit of a Triple

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## Lemma

*All elements of the triples in the orbit of  $(x, y, z)$  with respect to  $x$  are of the form  $h + \frac{\sigma}{h}$  where  $h + h^{-1} = x$  and  $\sigma \neq 1$  (BGS)*

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We are in *The End Game* when a triple has order greater than  $p^{1/2+\delta}$  for  $\delta > 0$ .

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## Theorem

*All triples in the end game have a triple of maximal order in their orbits (BGS).*

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We are in *The Middle Game* when a triple has order less than  $p^{1/2+\delta}$ .



# An Important Result

## Theorem

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So, we do not have to show every triple is connected to the cage, only that all but at most  $p - 1$  are.

Our goal is to show for what values of  $p$  we can guarantee this for.

## Connecting to Triple of Larger Order

Let a triple  $(x_1, x_2, x_3)$  have order  $n < p^{1/2+\delta}$  with  $x_1$  being the coordinate of largest order.

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### Example

If a triple  $(x_1, x_2, x_3)$  has order 12 and we know there are 10 triples in the orbit of order at most 12, there must be a triple of order greater than 12.

# What to Bound

## Theorem

Let  $x_1 = h_1 + h_1^{-1}$  with the order of  $x_1$  being  $n$ . The number of elements in the orbit of  $(x_1, x_2, x_3)$  of the form  $h_2 + h_2^{-1}$  where  $\text{ord}(h_2) = m \leq n$  is at most

$$20 \cdot \max\left\{(n \cdot m)^{\frac{1}{3}}, \frac{(n \cdot m)}{p}\right\}$$

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Since  $n$  and  $m$  in the middle game are less than  $p^{1/2+\delta}$  the second part is irrelevant so we get an upper bound of  $20 \cdot (n \cdot m)^{\frac{1}{3}}$



## Summation we are After

To show the triple of order  $y$  as an element of larger order in its orbit, it suffices to show

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Unfortunately, this doesn't always hold, so we want to find the smallest value of  $n$  where it fails. This will depend on the prime  $p$ .

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$$20 \cdot n^{\frac{1}{3}} \cdot \sum_{\substack{d \leq n \\ d | p \pm 1}} d^{\frac{1}{3}} \leq 20 \cdot n^{\frac{2}{3}} \sum_{\substack{d \leq n \\ d | p \pm 1}} 1 \leq 20 \cdot n^{\frac{2}{3}} \cdot \tau(p \pm 1) \leq 20 \cdot C \cdot p^{\frac{1}{\log \log(p)}} \cdot n^{\frac{2}{3}}$$

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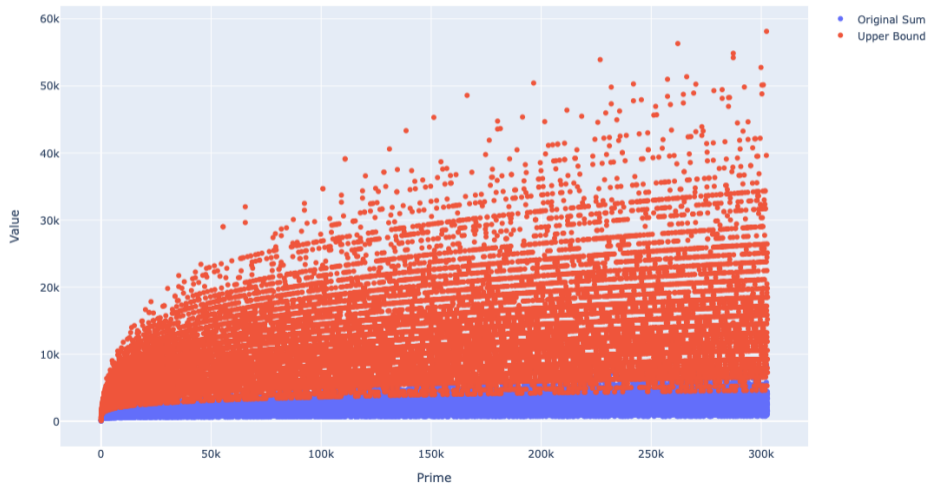
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Notice if we just had the order of a triple divides  $p^2 - 1$  instead of  $p \pm 1$  we would have roughly the square of this for our answer.

# Graph for Upper Bound

Original Sum and Upper Bound for the First Few Primes



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This will allow us to bound the number of triples not connected to the cage.

So, let  $y_p \leq (20 \cdot C \cdot p^{\frac{1}{\log \log(p)}})^3$  be the greatest divisor with this property.

## Another Summation

Let  $y_p \leq (20 \cdot C \cdot p^{\frac{1}{\log \log(p)}})^3$ . Triples with order less than  $y_p$  may not be connected to the cage.

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How many such triples are there? For the first coordinate, there are  $\sum_{\substack{d \leq y_p \\ d|p \pm 1}} d$  possibilities, since each divisor  $d$  contributes no more than  $d$  elements of that order.

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This is bounded by

$$\sum_{\substack{d \leq y_p \\ d | p \pm 1}} d < C \cdot y_p \cdot \tau(p \pm 1) < C' \cdot p^{\frac{4}{\log \log(p)}}$$

where  $C'$  is a "small" constant.

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To see this, set  $x$  and  $y$  in the Markoff equation  $x^2 + y^2 + z^2 - xyz = 0$ . This becomes a quadratic equation in  $z$  and we see the solutions are

$$z = \frac{xy \pm \sqrt{(xy)^2 - 4(x^2 + y^2)}}{2}.$$



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In total, there are at most  $2 \cdot (C' \cdot p^{\frac{4}{\log \log(p)}})^2 = C'' p^{\frac{8}{\log \log(p)}}$  triples where  $C''$  is again "small".

# The Big Reveal

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7787230284780299625150563291249403394823772620298436994917402090195995074  
7642789241468280077315444827633791559611419347950254097535856157166975009  
7158620540703858858904072561643063296397476748443544607407652354384777167  
1896405486844894947300377395071466792054518904692494753131475818530403786  
6541459144927618666489285293885227309626197614981891464741548360089233331  
1539909647185210552740723034041856213955035936355864103340086723634207005  
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8484001341201348391224378074935473507514716708458912061828816811578282646  
1837885486796686100883775234472717379929328796353942091918645011166241418  
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1559940407255733491561390514718861532568995894630353855313289779729964926  
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This, with the theorem of Chen from earlier, shows if  $p > e^{e^8} \approx 10^{1300}$  then  $G_p$  is connected.

Thank You!